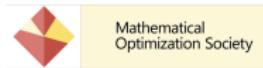


# Coluna.jl : An Open-Source Branch-Cut-and-Price Framework

C. Bentes, T. Bulhoes, H. Kramer, G. Marques, V. Nesello, A. Pessoa, M. Poss, A. Subramanian, M. Tanneau, R. Sadykov,

E. Uchoa, F. Vanderbeck

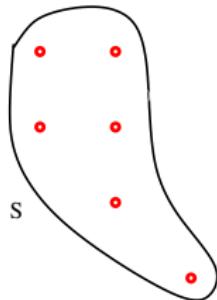


November 2019

## Combinatorial Optimization Problem

$$(CO) \equiv \min\{c(s) : s \in \mathbf{S}\}$$

where  $\mathbf{S}$  is the **discrete** set of feasible solutions.



# Integer Programming Formulation

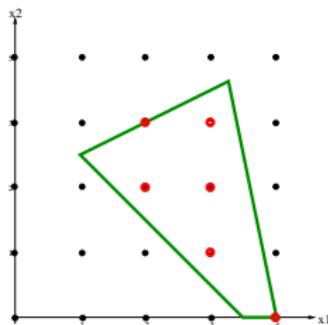
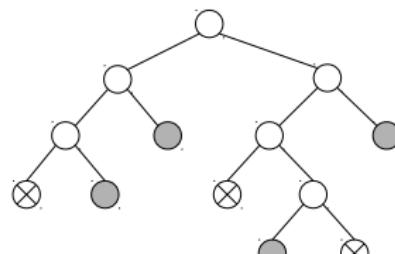
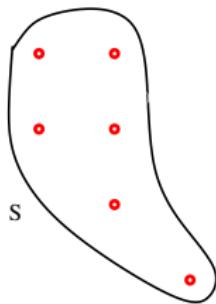
## Combinatorial Optimization Problem

$$(CO) \equiv \min\{c(s) : s \in \mathbf{S}\}$$

where  $\mathbf{S}$  is the **discrete** set of feasible solutions.

## Integer Programming Formulation

A **polyhedron**  $\mathbf{P} = \{x \in \mathbb{R}^n : Ax \geq \mathbf{a}\}$  is a formulation for  $(CO)$  iff  
 $\min\{c(s) : s \in \mathbf{S}\} \equiv \min\{cx : x \in \mathbf{P}_I = \mathbf{P} \cap \mathbb{N}^n\}.$



# Decomposition Approaches

- On a **subset of constraints** (Dantzig-Wolfe): **multi-resources**

$$\begin{array}{lllllll} \min & c^1 x^1 & + & c^2 x^2 & + & \dots & + & c^K x^K \\ & Ax^1 & + & Ax^2 & + & \dots & + & Ax^K & \geq a \\ & B^1 x^1 & & & & & & & \geq b^1 \\ & & B^2 x^2 & & & & & & \geq b^2 \\ & & & & & \dots & & & \geq \dots \\ & & & & & & & B^K x^K & \geq b^K \end{array}$$

- On a **subset of variables** (Benders): **multi-decision-levels**

$$\begin{array}{lllllll} \min & a \textcolor{blue}{x} & + & b^1 \textcolor{green}{y}^1 & + & b^2 \textcolor{green}{y}^2 & + & \dots & + & b^K \textcolor{green}{y}^K \\ & A \textcolor{blue}{x} & + & B^1 \textcolor{green}{y}^1 & & & & & & \geq c^1 \\ & A \textcolor{blue}{x} & + & & B^2 \textcolor{green}{y}^2 & & & & & \geq c^2 \\ & A \textcolor{blue}{x} & + & & & & \dots & & & \geq \dots \\ & A \textcolor{blue}{x} & + & & & & & B^K \textcolor{green}{y}^K & \geq c^K \end{array}$$

Assume a bounded integer integer problem with structure:

$$\begin{aligned} [F] \equiv \min \quad & c x : \\ & Ax \geq a \\ x \in X = \{ & Bx \geq b \\ & x \in \mathbb{N}^n \} \end{aligned}$$

Assume that **subproblem**

$$[SP] \equiv \min \{c x : x \in X\} \quad (1)$$

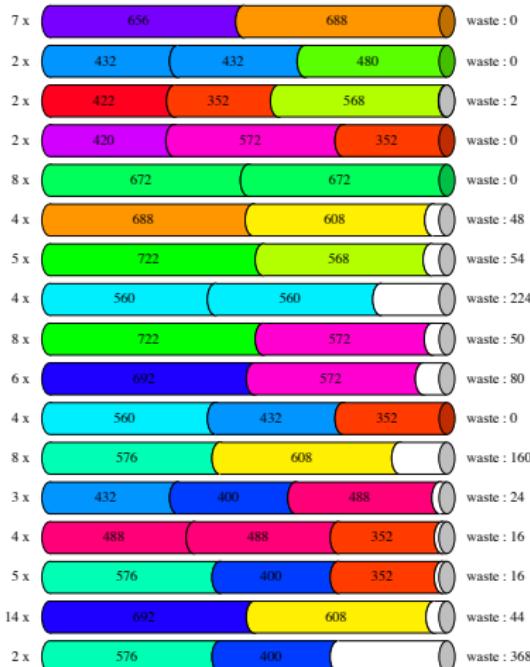
is “relatively easy” to solve compared to problem [F]. Then,

$$X = \{x^q\}_{q \in Q}$$

$$\text{conv}(X) = \{x \in \mathbb{R}_+^n : x = \sum_{q \in Q} x^q \lambda_q, \sum_{q \in Q} \lambda_q = 1, \lambda_q \geq 0, q \in Q\}$$

$$[F] \equiv \min \left\{ \sum_{q \in Q} c x^q \lambda_q : \sum_{q \in Q} A x^q \lambda_q \geq a, \sum_{q \in Q} \lambda_q = 1, \sum_q x^q \lambda_q \in \mathbb{N}^n \right\}$$

# The Cutting Stock Problem (CSP)



Number of distinct cutting patterns : 17

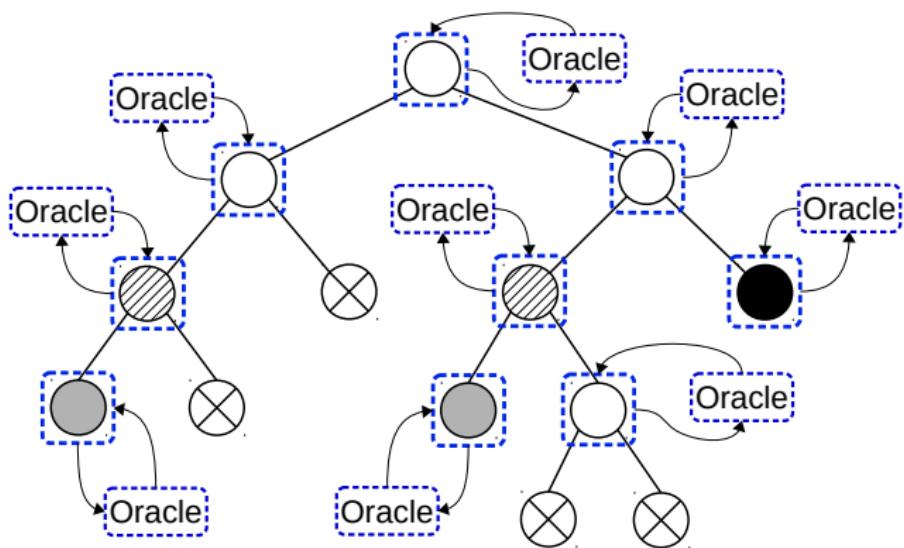
Number of rolls : 88

Total waste : 5090

$$\begin{aligned}
 & \min \quad \sum_{k=1}^K y_k \\
 \text{s.t.} \quad & \sum_{k=1}^K x_{ik} \geq d_i \quad \forall i \\
 & \sum_i w_i x_{ik} \leq W y_k \quad \forall k \\
 & x_{ik} \in \mathbb{N} \quad \forall i, k \\
 & y_k \in \{0, 1\} \quad \forall k
 \end{aligned}$$

$$\begin{aligned}
 & \min \sum_q \lambda_q \\
 & \sum_i x_i^q \lambda_q \geq d_i \quad i = 1, \dots, n \\
 & \lambda_q \geq 0 \quad \forall q
 \end{aligned}$$

# Solving decomposed problems



# Benders Decomposition



$$\begin{aligned} \max \quad & \textcolor{blue}{a} \textcolor{blue}{x} + \textcolor{green}{b^1} \textcolor{green}{y^1} + \textcolor{green}{b^2} \textcolor{green}{y^2} + \dots + \textcolor{green}{b^K} \textcolor{green}{y^K} \\ & \textcolor{blue}{A} \textcolor{blue}{x} + \textcolor{green}{B^1} \textcolor{green}{y^1} \\ & \textcolor{blue}{A} \textcolor{blue}{x} + \textcolor{green}{B^2} \textcolor{green}{y^2} \\ & \textcolor{blue}{A} \textcolor{blue}{x} + \dots \\ & \textcolor{blue}{A} \textcolor{blue}{x} + \textcolor{green}{B^K} \textcolor{green}{y^K} \end{aligned}$$

$$\begin{aligned} \max \quad & \{ \textcolor{blue}{a} \textcolor{blue}{x} + \textcolor{green}{b} \textcolor{green}{y} : \\ & \textcolor{blue}{A} \textcolor{blue}{x} + \textcolor{green}{B} \textcolor{green}{y} \leq \textcolor{black}{c}, \\ & \textcolor{blue}{x} \in \mathbb{N}^n, \textcolor{green}{y} \in \mathbb{R}_+^p \}, \end{aligned}$$

$$\begin{aligned} \max \{ & a \mathbf{x} + b \mathbf{y} : \\ & A \mathbf{x} + B \mathbf{y} \leq c, \\ & \mathbf{x} \in \mathbb{N}^n, \mathbf{y} \in \mathbb{R}_+^p \}, \quad \phi(\mathbf{x}) := \max_{\mathbf{y} \in \mathbb{R}_+^p} \{ b \mathbf{y} : B \mathbf{y} \leq c - A \mathbf{x} \} \end{aligned}$$

First stage:

$$\begin{aligned} \max_{\mathbf{x}, \eta} \{ & \eta : \\ & \eta \leq a \mathbf{x} + \phi(\mathbf{x}), \Leftrightarrow u_0^q (\eta - a \mathbf{x}) + u^q (A \mathbf{x} - c) \leq 0 \forall q \\ & \mathbf{x} \in \mathbb{N}^n, \eta \in \mathbb{R} \}, \end{aligned}$$

Second stage:

$$\begin{aligned} \min \mathbf{v} \leq 0 ? \\ b \mathbf{y} + \mathbf{v} \geq (\eta - a \mathbf{x}) \\ -B \mathbf{y} + 1 \mathbf{v} \geq (A \mathbf{x} - c) \\ (\mathbf{y}, \mathbf{v}) \in \mathbb{R}_+^{p+1} \end{aligned} \quad \begin{aligned} \max u_0 (\eta - a \mathbf{x}) + u (A \mathbf{x} - c) \leq 0 ? \\ u_0 b - u B \leq 0 \\ u_0 + 1 x \leq 1 \\ (u_0, u) \in \mathbb{R}_+^{1+m_2} \end{aligned}$$

$\{u^1, \dots, u^T\}$  are its **extreme dual solutions**.

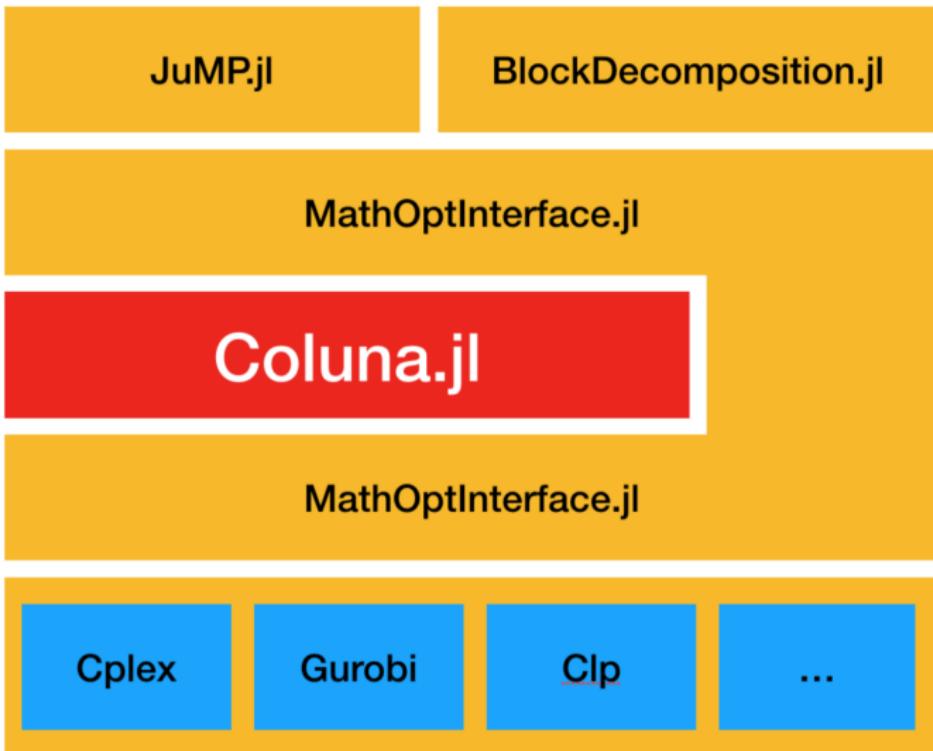
## Integer Programming Decomposition & Reformulations

- A powerful **way** to **exploit** the combinatorial structure.
- **Benders & Dantzig Wolfe decomposition** are **generic** schemes to derive & handle **quality** reformulations
- **Size** can be coped with using **dynamic generation**: a **small %** of **variables** and **constraints** are needed; hence it **scales up** to real-life applications.
- With **efficiency enhancement** features, these approaches can be highly **competitive**
- Implementation can be **generic**, with tools such as **Coluna.jl**

# Coluna.jl 0.2

- **What:**
  - a **Column-and-Row** generation code
  - Open-Source: [Mozilla Public Licence - MPL](#)
  - in [Julia 1.2](#)
  - Builds on [JuMP-JuliaOpt](#) extending it with [BlockDecomposition.jl](#)
- **Uses:**
  - [Dantzig-Wolfe & Benders decomposition](#)
  - [Extended Formulation Approaches](#)
  - [Robust & Stochastic Optimization](#)
  - [Machine Learning](#)
- **Support:** [Math Optimization Society \(MOS\)](#) & [JuliaOpt](#)
- **URL**  
<https://github.com/atoptima/Coluna.jl>

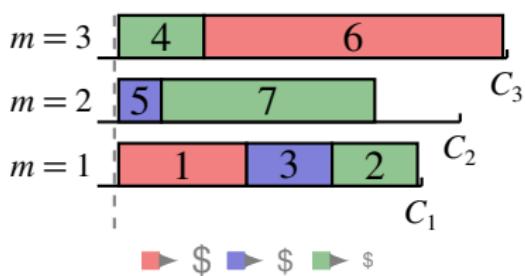
# Coluna Environment



## OUTLINE

- **Modeling**
- **Data Structure**
- **Algorithms**
- **Current Features**
- **Features to come**
- **Installation**

## Exemple : Generalized Assignment



**Machines** = 1:3

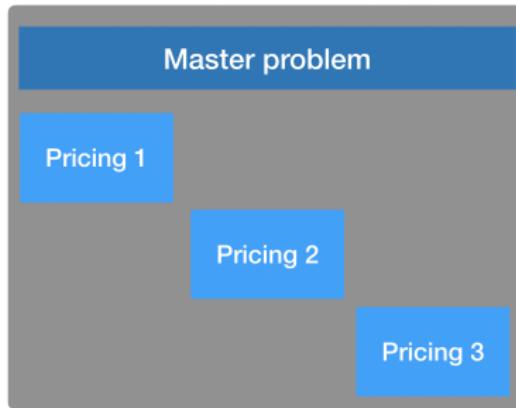
```
@axis(Machines)
```

```
model = BlockModel()
```

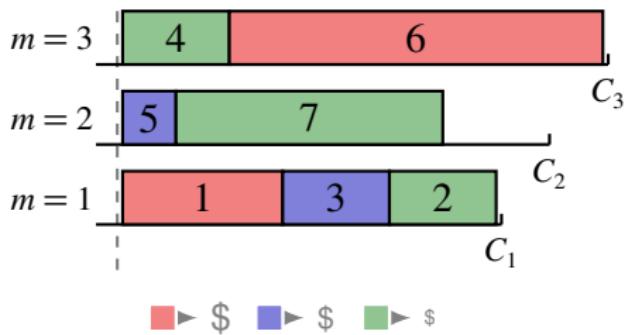
```
@variable(model, x[ j in Jobs, m in Machines], Bin)
@objective(model, Min,
    sum(cost[j,m] * x[j,m] for j in Jobs , m in Machines,))
@constraint(model, cov[j in Jobs],
    sum(x[j,m] for m in Machines) >= 1 )
@constraint(model, knp[m in Machines],
    sum(weight[j,m] * x[j,m] for j in Jobs) <= Capacity[m]) )
```

```
@DantzigWolfe_decomposition(model, dec, Machines)
```

- Write a model of your problem using JuMP
- Use variable and constraint indices to define the decomposition
- Store extra information in the indices -> `axis` macro
- Create a decomposition tree ->  
`dantzig_wolfe_decomposition` macro



# Generalized Assignment : Automated Reformulation



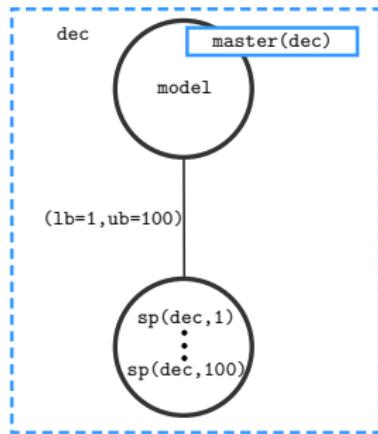
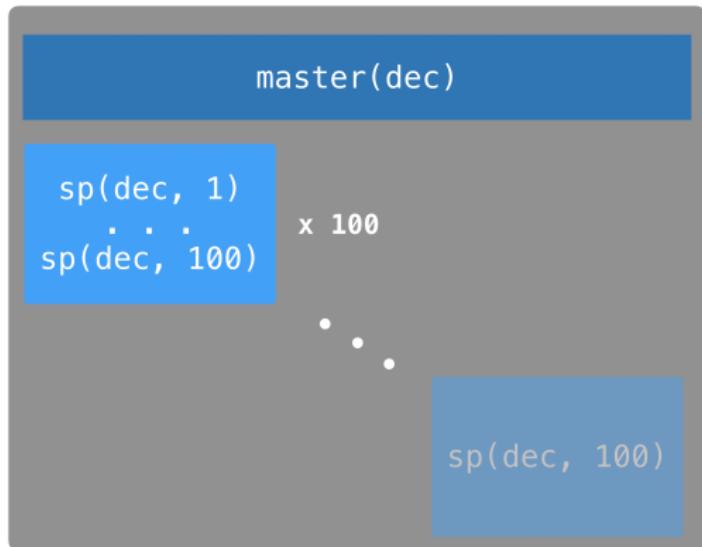
$$\begin{aligned}
 & \min \quad \sum_{jm} c_{jm} x_{jm} \\
 \text{s.t.} \quad & \sum_{m=1}^M x_{jm} \geq 1 \quad \forall j \\
 & \sum_j w_{jm} x_{jm} \leq C_m \quad \forall m \\
 & x_{jm} \in \{0, 1\} \quad \forall j, m
 \end{aligned}$$


---

$$\begin{aligned}
 & \min \sum_q c_q \lambda_q \\
 & \sum_{m \in M} \sum_{q \in Q^m} x_j^q \lambda_q \geq 1 \quad \forall j \\
 & \sum_{q \in Q^m} \mathbf{y}^q \lambda_q = 1 \quad \forall m \\
 & \lambda_q \in \{0, 1\} \quad \forall q
 \end{aligned}$$

where  $\mathbf{y}$  is an added **setup-variable** in the sub-problem.

# Cutting Stock : Identical subproblems



# Cutting Stock : Identical subproblems



Number of distinct cutting patterns : 17

Number of rolls : 88

Total waste : 5090

**SheetTypes** = 1:1

Orders = 1:100

`model = BlockModel()`

`@axis(SheetTypes)`

`@variable(model, x[i in Orders, k in SheetTypes], Int)`

`@variable(model, y[k in SheetTypes], Int)`

`@objective(model, Min,`

`sum(cost[k] * y[k] for k in SheetTypes))`

`@constraint(model, cov[i in Orders],`

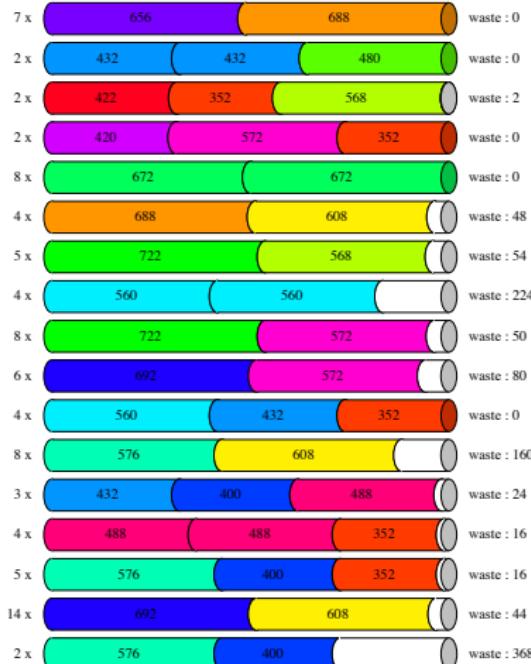
`sum(x[i,k] for k in SheetTypes) >= 1 )`

`@constraint(model, knp[k in SheetTypes],`

`sum(weight[i] * x[i,k] for i in Orders) <= Capacity[k]) )`

`@DantzigWolfe_decomposition(model, dec, SheetTypes)`

# Cutting Stock : how to Specify & to Get the solution



Number of distinct cutting patterns : 17

Number of rolls : 88

Total waste : 5090

master = getmaster(dec)

addbranchingpriority!(master, y, 1)

SP = getsubproblems(dec)

specify!(SP, lm = 0, um = U, **setup\_var** = y)

optimize!(model)

for k in SheetTypes

solutions = Coluna.getsolutions(model, SP[k])

for sol in solutions

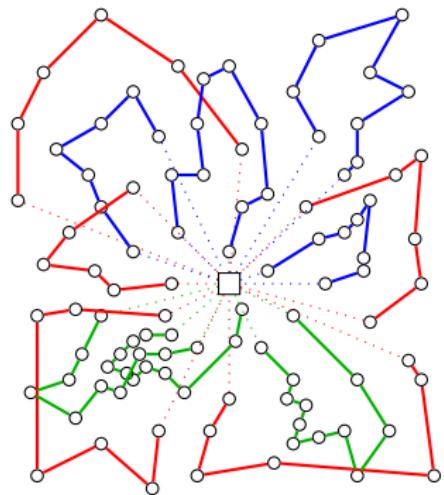
**multiplicity** = getvalue(sol, y[k])

**x\_val** = [getvalue(sol, x[i, k]) for i in Orders]

end

end

# CVRP : Expression & ad-hoc pricing solver

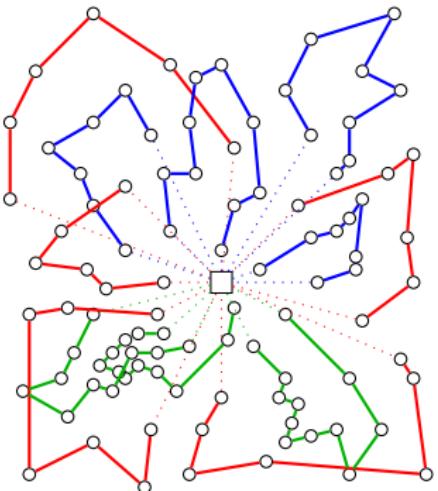


$$\begin{aligned}
 \min \quad & \sum_{e,r} c_e x_e^r \\
 \text{s.t.} \quad & \sum_{e \in \delta(i), r} x_e^r \geq 2 \quad \forall i \\
 & x^r \in X^r \quad \forall r
 \end{aligned}$$


---

$$\begin{aligned}
 \min \quad & \sum_{e,r,q \in Q^r} c_e x_e^q \lambda_q \\
 \sum_{e \in \delta(i), r, q \in Q^r} x_e^q \lambda_q & \geq 2 \quad \forall i \\
 \sum_{q \in Q^r} \lambda_q & \leq U^r \quad \forall r \\
 \lambda_q & \in \{0, 1\} \quad \forall r, q \in Q^r
 \end{aligned}$$

# CVRP : Expression & ad-hoc pricing oracle



```
model = BlockModel()  
@axis(VehicleTypes)  
@variable(model, x[e in Edge(G)], Int)  
@variable(model, y[k in VehicleTypes], Bin)  
@variable(model, z[a in Arc(G), k in VehicleTypes], Bin)  
@objective(model, Min,  
    sum(cost[e] * x[e] for e in Edge(G)))  
@constraint(model, cov[i in Locations],  
    sum(x[e] for e in Edge(G,i)) = 2 )  
@expression(model, x[ (i,j) in Edge(G)],  
    sum((z[(i,j),k] + z[(j,i),k]) for k in VehicleTypes) )  
@DantzigWolfe_decomposition(model, dec, VehicleTypes)  
  
function rcsp_fct(cb)  
    k = getspidx(cb)  
    red_costs = [getredcost(cb, x[e]) for e in Edge(G)]  
    route = rcsp_pricing_solver(data, k, red_costs)  
end  
@oracle(model, sp_cb[k in VehicleTypes], rcsp_fct)
```

# 2D Cutting Stock : Original & master reformulations

$$\min \sum_{s \in S} z_s$$

$$\text{s.t. } \sum_{s \in S, p \in P} x_{jps} \geq 1 \quad \forall j$$

$$\sum_p w_{ps} \leq W z_s \quad \forall s$$

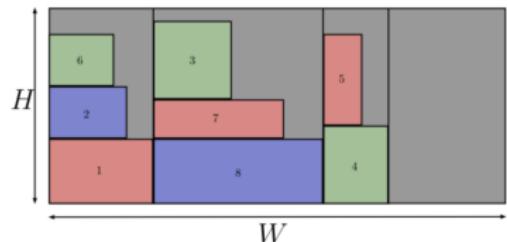
$$w_{ps} \geq w_j x_{jps} \quad \forall j, p, s$$

$$\sum_j h_j x_{jps} \leq H y_{ps} \quad \forall p, s$$

$$x_{jps} \in \{0, 1\} \quad \forall j, p, s$$

$$y_{ps} \in \{0, 1\} \quad \forall p, s$$

$$z_s \in \{0, 1\} \quad \forall s$$



$$\min \sum_s \lambda_s$$

$$\sum_s x_j^s \lambda_s \geq 1 \quad \forall j$$

$$\lambda_s \in \mathbb{N} \quad \forall s$$

# 2D Cutting Stock : Pricing problem and its reformulation

$$\min 1 - \left( \sum_j \pi_j (\sum_p x_{jp}) \right)$$

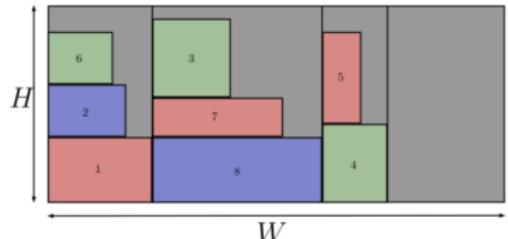
$$\text{s.t. } \sum_p x_{jp} \leq 1 \quad \forall j$$

$$\sum_p w_p \leq W \quad \forall$$

$$w_p \geq w_j x_{jp} \quad \forall j, p$$

$$\sum_j h_j x_{jp} \leq H \quad \forall p$$

$$x_{jp} \in \{0, 1\} \quad \forall j, p$$



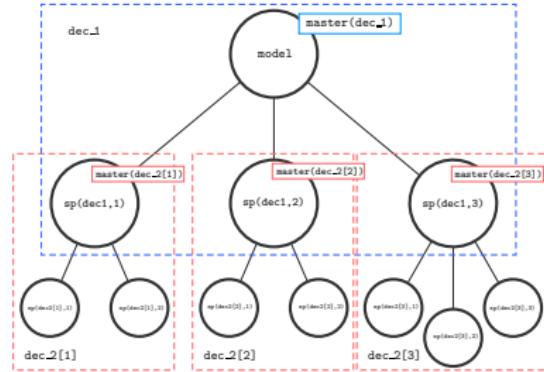
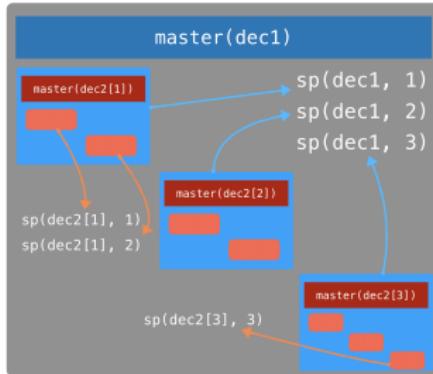
$$\max \sum_p \left( \sum_j \pi_j x_j^p \right) \lambda_p$$

$$\sum_p x_j^p \lambda_p \leq 1 \quad \forall j$$

$$\sum_p w_p \lambda_p \leq W$$

$$\lambda_p \in \mathbb{N} \quad \forall p$$

# 2D Cutting Stock : Nested Dantzig-W decomposition



@axis(**SheetTypes**)

@axis(**StripPatterns**)

@DantzigWolfe\_decomposition(*model*, *dec1*, **SheetTypes**)

@DantzigWolfe\_decomposition(*dec1*, *dec2[k]* in **SheetTypes**, **StripPatterns**)

1th stage facility Linking C.

1th stage C. for each facility

link C. for each facility

?2d stage facility C.

$$\begin{array}{llll}
 \min \sum_k c^k \mathbf{x}^k & + \sum_{k,\omega} h_\omega^k \mathbf{y}_\omega^k \\
 \sum_k \mathbf{x}_t^k & \geq N_t \quad \forall t \\
 A^k \mathbf{x}^k & \leq a^k \quad \forall k \\
 B^k \mathbf{x}^k & + T^k \mathbf{y}_\omega^k \leq b^k \quad \forall k, \omega \\
 \mathbf{x}^k & \sum_k \mathbf{y}_{\omega t}^k \geq d_{\omega t} \quad \forall \omega, t \\
 & \mathbf{y}_\omega^k \in \{0, 1\}^n \times \mathbb{R}_+^m \quad \forall k, \omega
 \end{array}$$



@axis(**K-facilities**)

@axis(**Ω-scenarios**)

@DantzigWolfe\_decomposition

(model, dec1, **K-facilities**)

@Benders\_decomposition

(model, dec2, **Ω-scenarios**)

```
struct Problem <: AbstractProblem
    original_formulation::Union{Nothing, Formulation}
    re_formulation::Union{Nothing, Reformulation}
end

struct Formulation{Duty <: AbstractFormDuty} <: AbstractFormulation
    parent::Union{AbstractFormulation, Nothing}
    manager::FormulationManager
end

struct Reformulation <: AbstractFormulation
    parent::Union{Nothing, AbstractFormulation}
    master::Union{Nothing, Formulation}
    dw_pricing_subprs::Vector{AbstractFormulation}
    benders_sep_subprs::Vector{AbstractFormulation}
end
```

```
struct FormulationManager
    vars::VarDict
    constrs::ConstrDict
    coefficients::VarConstrMatrix
    expressions::VarVarMatrix
    primal_sols::VarVarMatrix
    dual_sols::ConstrConstrMatrix
end

struct MembersVector{IdT,ElemT,ValueT} <: AbstractMembersContainer
    elements
    records
end

struct MembersMatrix{IdT,ElemT,ValueT} <: AbstractMembersContainer
    cols
    rows
end
```

# Coluna's Data Structure

```
struct Variable <: AbstractVarConstr
    id::Id[Variable]
    name::String
    duty::Type{<: AbstractVarDuty}
    perene_data::VarData
    cur_data::VarData
    moirecord::MoiVarRecord
end
```

```
struct VarData <: AbstractVarData
    cost::Float64
    lb::Float64
    ub::Float64
    kind::VarKind
    sense::VarSense
    inc_val::Float64
    is_active::Bool
    is_explicit::Bool
end
```

```
struct Constraint <: AbstractVarConstr
    id::Id[Constraint]
    name::String
    duty::Type{<: AbstractConstrDuty}
    perene_data::ConstrData
    cur_data::ConstrData
    moirecord::MoiConstrRecord
end
```

```
struct OriginalVar <: AbstractOriginalVar end
struct OriginalExpression <: AbstractOriginalVar end

struct MasterPureVar <: AbstractOriginMasterVar end
struct MasterCol <: AbstractAddedMasterVar end
struct MasterArtVar <: AbstractAddedMasterVar end
struct MasterBendSecondStageCostVar <: AbstractAddedMasterVar end
struct MasterBendFirstStageVar <: AbstractOriginMasterVar end
struct MasterRepPricingVar <: AbstractMasterRepDwSpVar end
struct MasterRepPricingSetupVar <: AbstractMasterRepDwSpVar end

struct DwSpPricingVar <: AbstractDwSpVar end
struct DwSpSetupVar <: AbstractDwSpVar end
struct DwSpPureVar <: AbstractDwSpVar end

struct BendSpSepVar <: AbstractBendSpVar end
struct BendSpPureVar <: AbstractBendSpVar end
struct BendSpSlackFirstStageVar <: AbstractBendSpSlackMastVar end
struct BendSpSlackSecondStageCostVar <: AbstractBendSpSlackMastVar end

struct OriginalConstr <: OriginalConstr end

struct MasterPureConstr <: AbstractMasterOriginConstr end
struct MasterMixedConstr <: AbstractMasterOriginConstr end
struct MasterConvexityConstr <: AbstractMasterAddedConstr end

struct DwSpPureConstr <: bstractDwSpConstr end
struct DwSpRepMastBranchConstr <: bstractDwSpConstr end

:
```

**Algorithm:** is a routine involved in the optimisation process.

**Solver:** is a complete algorithmic strategy that we can solve a (re)formulation.

```
struct abstractConquerAlgorithm <: AbstractAlgorithm end
struct abstractDivideAlgorithm <: AbstractAlgorithm end
struct abstractExploreAlgorithm <: AbstractAlgorithm end

struct MySolver <: AbstractAlgorithm
    conquer_strategy::abstractConquerAlgorithm
    divide_strategy::abstractDivideAlgorithm
    explore_strategy::abstractExploreAlgorithm
end
```

## Current

- **Modeling**

- JuMP Extension
- Dantzig Decomposition
- Benders Decomposition

- **Algorithms**

- Column Generation
- Benders cut Gen.
- Branch-and-Bound
- Branch-and-Price
- Benders-Decomposition
- Preprocessing
- Primal Heuristics

## Current

- **Modeling**
  - JuMP Extension
  - Dantzig Decomposition
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  - Column Generation
  - Benders cut Gen.
  - Branch-and-Bound
  - Branch-and-Price
  - Benders-Decomposition
  - Preprocessing
  - Primal Heuristics

## Coming Soon

- **Efficiency**
  - Stabilization
  - Interior Point Approach
  - Strong Branching
  - Further Primal Heuristics
  - Pricing Strategies
  - **Parallelisation**
- **Cutomization**
  - Pricing Callback
  - Cutting Callback
  - Branching Callback
- **Functionality**
  - Template Cut Generation
  - Ad-hoc Branching
  - Nested / Hybrid Decomp.

- Install **Julia 1.2+**.
- Install **JuMP & BlockDecomposition**
- Install **GLPK** (or **CPLEX**, or **Gurobi**, or **Xpress** or ...) as the default underlying MOI Optimizer for the master and the subproblems.
- Get **Coluna.jl** (using the package manager of Julia)
  - Go to the **Pkg-REPL-mode** :  
the Pkg REPL-mode is entered from the Julia REPL using the key ]
  - The command to **install Coluna.jl** and its dependencies is  
**pkg> add https://github.com/atoptima/Coluna.jl.git**
  - Start **using Coluna** by doing :  
**using Coluna**

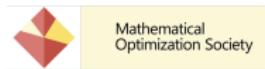
# Welcome

- **$\beta$ -users**
- **Integration**
  - Algorithms: Subgradient, Interior Point, Primal Heuristics
  - Oracles: Pricing, Separation
  - Applications: demo problems
- **Code review**
- **Co-development**

# Coluna.jl : An Open-Source Branch-Cut-and-Price Framework

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